Scalability Update, 2007

Kris Kennaway The FreeBSD Project kris@FreeBSD.org

May 17, 2007

Progress on FreeBSD SMP performance and scalability since BSDCan Dev Summit 2006

- $1.\ {\sf A}$ case study: SQL database performance
- 2. Major SMP changes since May 2006
- 3. Remnants of the Giant lock
- 4. Future work and new SMP goals

Why Database?

- Off-the-shelf benchmark (sysbench; /usr/ports/benchmarks/sysbench/)
- Seems to be a reasonable benchmark (as opposed to super-smack; 1-byte I/O!)
- FreeBSD did not perform well compared to Linux; excellent motivator for performance improvements

What Database?

- mySQL 5.0.37 (thread-based)
- PostgreSQL 8.2.4 (process-based + SysV IPC)

- amd64 system, 4*dual-core 2.2GHz CPUs
- 3GB RAM (working set fits in RAM)
- sysbench OLTP test
 - Complex transaction-based queries
 - Read-only; no disk access to avoid benchmarking disk performance
- multithreaded benchmark client
- clients and servers on the same system
- communication via UNIX domain socket
- 2 minute runs, after a warm-up run with 8 clients.

MySQL: Progress in FreeBSD 5.x and 6.x



E.

< 🗗 🕨

MySQL: The competition (FreeBSD vs Linux)



▲□→ < □→</p>

< ∃⇒

MySQL: Progress in FreeBSD 7.0 through February 2007



- ◆ @ ▶ - ◆ 臣

э



・ロト ・回ト ・ヨト

< ≣⇒

Linux: Blame it on malloc? glibc vs tcmalloc



Source: http://ozlabs.org/~anton/linux/sysbench/

A (1) > A (2)

- Some reports on LKML of difficulty reproducing the tcmalloc result on 8-core systems
- tcmalloc claimed not to be a good general-purpose allocator
- A process-global semaphore acquired by all linux futex operations.
 - Agrees with the asymptotically single-threaded performance measured on amd64

...not our problem!



@ ▶

э

Kris Kennaway The FreeBSD Project kris@FreeBSD.org Scalability Update, 2007



E

・ロト ・回ト ・ヨト

- More efficient lock profiling
 - Still very inefficient; explore a worker thread model?
- UNIX domain socket locking; fine-grained locking for better concurrency
- Experimental scheduler locking work (on-going)
- reader-writer (rw) locks (non-sleepable)
- Optimized sx locks (sleepable); now an efficient low-level primitive
- \blacktriangleright File descriptor locking; home-brewed msleep lock \rightarrow sx

- ULE 2.0. Now significantly out-performing 4BSD on the SMP workloads that have been benchmarked.
 - A remaining performance anomaly: poor CPU affinity at intermediate loads on dual-core systems
 - Benchmarking on other workloads required
- Giant pushdown
- Network stack optimization and regularization
- libthr is now the default thread library in 7.0 (finally!)
 - Should also be default on 6.x
- CAM subsystem locking (scottl)

The SMPng project was launched in June 2000 and was a major focus of development in FreeBSD 5.x and above.

"The goal of the SMPng Project is to decompose the Giant lock into a number of smaller locks, resulting in reduced contention (and improved SMP performance)."

The SMPng project was formally concluded in February 2007 after achieving this goal.

The FreeBSD 7.0 kernel is mostly Giant-free. Remaining Giant-locked systems are:

- TTY locking. In most configurations the Giant lock has now effectively devolved onto a Giant TTY lock.
- Some drivers (some CAM SIMs, pseudo-devices)
- Some filesystems (MSDOS, SMB, NFSv4, ...)
- Some network protocols (IPX over IP, IPv6 ND6/MLD6, netatm, ...)
- newbus (scottl; work in progress)
- USB, Firewire
- Parts of VM (contigmalloc, ...)
- Some scattered use elsewhere: sysctls, file operations, sysarch(), ...

See http://wiki.freebsd.org/SMPTODO for the up-to-date list of Giant-locked code and project ownership.



æ

Kris Kennaway The FreeBSD Project kris@FreeBSD.org Scalability Update, 2007

- KAME IPSec; in favour of Fast IPSec (gnn, bz)
- I4B ISDN stack (rwatson)
- Legacy drivers: an arl awi cnw ce cp cs ctau cx en ex fe hfa idt ie if_ic oltr mn pcf pdq ray sbni sbsh snc sr tx wl xe (rwatson)
- > netatm? (rwatson, Skip Ford)

Legacy/home-grown locking primitives

- lockmgr
 - Proof of concept from ups
 - ► 1000 concurrent stat calls of the same file on an 8-core: reduces system time by 95% and real time by 82%.
 - lockmgr to be rewritten by Attilio Rao for Google SoC 2007
- Any others lurking?
 - msleep() is a pessimal synchronization/serialization primitive!
 - Lesson of SMPng: when the lock owner is running on another CPU it is almost always much better to spin instead of always going to sleep immediately.

回 と く ヨ と く ヨ と

- Use standard primitives instead of rolling your own.
- Consolidation of primitives; too many?
 - sema()
 - msleep()/tsleep()/wakeup()/... \rightarrow cv_*()?
 - ▶ ...?

Future work II: General

- Look for opportunities for shared locking but profile carefully!
 - e.g. namei() is an "obvious" candidate for using a rwlock, but this slows it down significantly.
 - probably because it exposes lockmgr to concurrent access.
- select() locking (giant select lock)
 - How can it be optimized? This is the major remaining source of lock contention/wait time for the SQL benchmarks (contended on 50-75% of acquisitions)

- SysV IPC; POSIX advisory locks
 - Sanitize, explore finer-grained locking
- Remove non-MPSAFE subsystem crutches
- ► VM
 - vm_page_queue mutex is heavily contended

Filesystem work

- More MPSAFE filesystems needed
 - MSDOS: Brian Chu, Google SoC 2007
- Shared lookups for UFS
 - Enables other work: rwlock for namei()

Scheduler

- scheduler lock is abused for non-scheduler purposes (rusage, ldt, vmmeter, timers)
 - //depot/user/attilio/attilio_schedlock
- ▶ global scheduler lock appears to be a barrier to scaling on ≥ 8 CPUs.
 - WIP to explore per-CPU scheduler locks with ULE; jeffr, attilio

- Explore methods for improving parallelism in network processing
- Driver work
 - e.g. bce driver performs poorly with concurrent send/receive
 - others not yet evaluated
- multiple netisr worker threads
 - e.g. loopback transport; very easy to saturate a single netisr
 - important for services that cannot use unix domain socket for local transport
- Optimization work to support 10ge (kmacy)

- 1. Identify additional workloads to be optimized.
 - You can help! We need to identify good benchmarks modelling real-world performance cases.
 - Ideally should be simple to set up and operate.
 - If you have a test case I will be happy to work with you to profile it.
- 2. Consolidate performance on 8-core systems without regressing performance on < 8 CPU cores
- 3. Ready to push scaling upwards to 16 and 32-core systems
 - ...but need access to hardware
 - sun4v? Not currently stable.
 - Need access to 16-core amd64 hardware

The scaling of FreeBSD n.x appears to be governed by

$$N_{CPUs} \sim 2^{n-4}, \quad n \ge 4$$

同 ト く ヨ ト く ヨ ト

Kris Kennaway The FreeBSD Project kris@FreeBSD.org Scalability Update, 2007